

The Gamut and Time Arrow of Automated Nurse Rostering

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ABSTRACT

There is an undeniable global shortage of skillful nurses. This is a problem of high priority, which is correlated to workforce management issues. These issues can be palliated by increasing nurses' satisfaction based on flexible rosters using automated nurse rostering. This paper is concerned with nurse rostering based on constraint programming by satisfying global constraints, such as REGULAR, which is powerful but requires devising automata of acceptable states. It was proven that REGULAR is reformulatable to SLIDE. En route to reformulate a REGULAR-based into SLIDE-based solution to nurse rostering, and through elaboration of Minizinc implementation, the author of this paper proposes a couple of metaphors namely the gamut and time arrow to model the nurse rostering. A gamut is a whole-subset of objects, such as within a given color-space or by an output device. The time arrow may refer to the direction of time as comprehended in physics. The paper in hand elucidates the new formulation based on these two metaphors, and presents implementation in Minizinc.

Keywords: *Automated nurse rostering, Constraint programming, REGULAR constraint, SLIDE constraint, Gamut, Time Arrow, Minizinc.*

1. INTRODUCTION

Personnel scheduling is a core task in workforce management ([1], [2]), in many service-providing work milieus such as call centers [3], hospitals [4], and the like. It comes with many flavors such as considering manpower skills [5], breaks [6], days-off [7] and overlapping [8].

There is an undeniable global shortage of skillful nurses [9]. This is a problem of high priority [10], which is correlated to workforce management issues ([11], [12]). These issues can be palliated by increasing nurses' satisfaction based on flexible rosters and strewing non-day shifts [13].

Automated Nurse Rostering[14] is the rescue. Since the earlier formulation of nurse rostering [15], and the employment of linear programming [16], a huge body of literature is flourished([17], [18]). Many algorithms and heuristics are proposed such as evolutionary ([19], [20]), mimetic[21], electromagnetic [22], scatter search [23],

branching strategies [24], neighborhood ([25], [26]), column generation preferences [27]. One particular approach in vogue is the local search [28] and its guided version using Tabu search [29]. Shift design, and its optimization [30], may be considered as a preliminary step for rostering [31]. You may refer to [32], [33] and [34] for recent comprehensive surveys on nurse rostering models and methodologies.

This paper is concerned with nurse rostering based on constraint programming [35] by satisfying global constraints [36], such as REGULAR [37] which is powerful but requires devising automata of acceptable states [38]. It was proven that REGULAR is reformulatable to SLIDE[39].

En route to reformulate a REGULAR-based into SLIDE-based solution to nurse rostering, and through elaboration of Minizinc implementation, the author of this paper proposes a couple of metaphors namely the gamut and time arrow to model the nurse rostering. A gamut is a whole-subset of objects, such as within a given color-space or by an output device [40]. The time arrow may refer to the direction of time as comprehended in physics [41].

The rest of the paper in hand is divided into the following sections. Section 2 elucidates the new formulation based on these two metaphors, while Section 3 presents implementation in Minizinc. Sections 4 and 5 are for conclusions and future work respectively.

2. PROBLEM FORMULATION

This section elucidates the new formulation based on gamut and time arrow metaphors.

2.1 Gamut

A gamut is a whole-subset of objects, such as colors within the RGB color-space [40]. It can be viewed as a SET of all possible values for a VECTOR, spanning certain basis dimensions. Henceforth, it may be denoted as \acute{z} .



Example 1. The gamut of the playing-cards, $C = \mathcal{C}^{N,S}$, is the set of all-possibilities of cards. A card has a number N and a suit S.

$$N = \{1, 2, \dots, 10, J, K, Q\}$$

$$S = \{\clubsuit, \diamonds, \heartsuit, \spadesuit\}$$

Note that this gamut can be decomposed, based on the second dimension, into 4 sub-gamuts namely C_{\clubsuit} , C_{\diamonds} , C_{\heartsuit} , and C_{\spadesuit} .

Example 2. The gamut, in nurse rostering, is a bit elusive, as it is constituted from three consecutive shifts and has an exclusion of the state of three consecutive night shifts.

So, nurse state gamut is

$$G = \mathcal{C}^{S_1, S_2, S_3} - \langle \text{night, night, night} \rangle$$

where $S_1 = S_2 = S_3 = \{\text{day, night, off}\}$

Fig.1 shows a full probability tree of shifts. A white node denotes a day shift, while a black node denotes a night shift. For a matter of simplification, we assume that each dimension has only two possibilities namely the night or day shift.

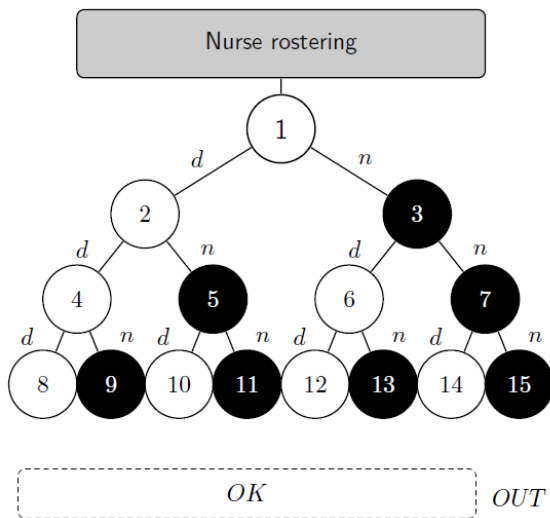


Fig. 1. Full probability tree of shifts.

Note that the gamut only contains the set of leaves except the last one that represents $\langle \text{night, night, night} \rangle$. Now, if we group and rename all the accepted nodes, as shown in Fig. 2, then group links then we reach to a Transition Diagram as shown in Fig. 3.

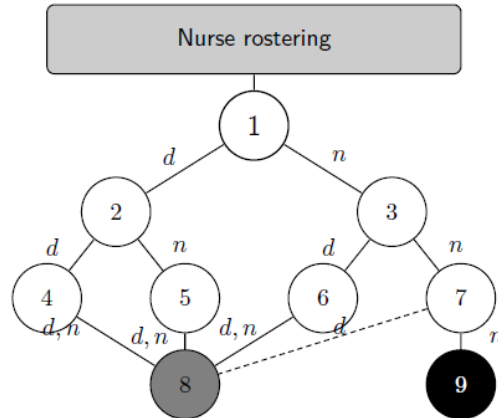


Fig. 2. Leaf acceptance.

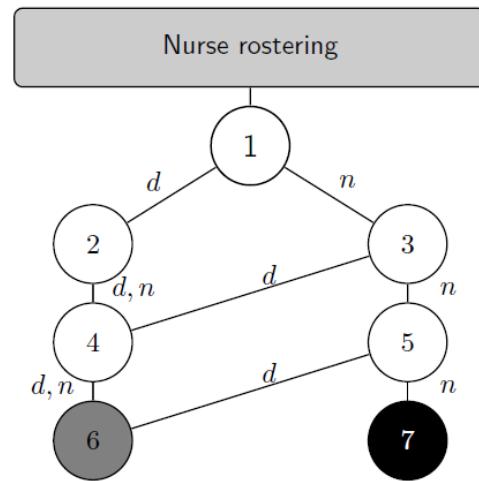


Fig. 3. Transition Diagram.

2.2 Time Arrow

SLIDE is useful if propagated in an efficient and effective way [42].

In encoding nurse rostering, dealing with time arrow \mathbb{F} can be one of three ways :

- retrospective :

$$\forall \square, 3 \leq \square \leq |\mathbb{F}| \quad |\mathbb{F}(\square-2) \cap \mathbb{F}(\square-1) \cap \mathbb{F}(\square)| = 0 \quad (1)$$

- knitting :

$$\forall \square, 2 \leq \square \leq |\mathbb{F}|-1 \quad |\mathbb{F}(\square-1) \cap \mathbb{F}(\square) \cap \mathbb{F}(\square+1)| = 0 \quad (2)$$

- anticipatory :

$$\forall \square, 1 \leq \square \leq |\mathbb{F}|-2 \quad |\mathbb{F}(\square) \cap \mathbb{F}(\square+1) \cap \mathbb{F}(\square+2)| = 0 \quad (3)$$

They are assumedly equivalent by shifting the index of the current time. The retrospective way compares the current day with the previous two days for ensuring that no nurse has occurred in a night shift. It only looks behind in a generate-and-evaluate fashion causing high probability to backtrack. The knitting way does the same but has a look-ahead to the next day. The anticipatory way has a look-ahead to two days to come.

3. IMPLEMENTATION

If the total number of nurses in a certain hospital is 7 nurses, along a time arrow spanning over 10 days, then time arrow of nurse availability is something like:

```
1..7
1..7
1..7
1..7
1..7
1..7
1..7
1..7
1..7
1..7
```

The MiniZinc[43] model to implement the retrospective way is shown in Listing 1

Listing 1. Retrospective implementation

```
1 include "globals.mzn";
2 int : num_nurses =7;
3 set of int: NURSES = 1.. num_nurses ;
4 int : num_days =10;
5 set of int : DAYS = 1.. num_days ;
6 array [DAYS] of var set of NURSES:
timeArrow;
7 constraint forall (i in 3.. num_days )(
8     disjoint ( timeArrow [i],
9         timeArrow [i-2] intersect timeArrow[i-
10         1]
11     ));
12 solve satisfy ;
13 output [show(timeArrow [i])++ "\n" | i in
DAYS];
```

The corresponding time arrow is

```
1..7
1..7
{}
1..7
7..7
1..6
1..7
```

```
{}
```

```
1..7
```

```
1..7
```

This means that if all the nurses are available for the first couple days to choose from. However, if one attends in both days in a night shift, she can not attend in the third day.

The problem can be decomposed into a dichotomy of two time arrows, one for day and another for night, as shown in Listing 2. This may help in parallelizing the execution.

Notice the introduction of constraints relating to the required nurses on each day and night shifts. Each of these constraints is applied on the corresponding time arrow. A constraint relating to the minimum night shifts per nurse is introduced too.

Listing 2. Two time arrows

```
1 include "globals.mzn";
2 int : total_nurses =7;
3 set of int : NURSES = 1.. total_nurses ;
4 int : total_days =10;
5 set of int : DAYS = 1.. total_days ;
6 int : required_days = 3;
7 int : required_nights = 2;
8 int : minimum_nights = 2;
9 array [ DAYS ] of var set of NURSES : timeArrow1 ;
10 array [ DAYS ] of var set of NURSES : timeArrow2 ;
11 constraint forall (i in 3.. total_days )(
12     disjoint ( timeArrow1 [i],
13         timeArrow1 [i -2] intersect timeArrow1 [i -1]));
14 constraint forall (i in 1.. total_days )(
15     card ( timeArrow1 [i ])== required_nights /\
16     card ( timeArrow2 [i ])== required_days /\
17     disjoint ( timeArrow1 [i], timeArrow2 [i]) );
18 constraint forall (j in 1.. total_nurses )(
19     sum (i in 1.. total_days )(
20         bool2int(j in
timeArrow1[i])>=minimum_nights);
21 solve satisfy ;
22 output [ show ( timeArrow1 [i]) ++ " | "
23     ++ show ( timeArrow2 [i]) ++ "\n"
24     | i in DAYS ];
```

Thus, the corresponding time arrows are:

```
1..2 | 3..5
1..2 | 3..5
{3,5} | {1,2,4}
{3,5} | {1,2,4}
{4,7} | 1..3
4..5 | 1..3
5..6 | 1..3
6..7 | 1..3
```



4..5 | 1..3
4..5 | 1..3

Generating the roster is shown in Listing 3, and the generated roster is shown in Fig. 4.

Listing 3. Generating the roster

```

1 include "globals.mzn";
2 int : total_nurses =7;
3 set of int : NURSES = 1.. total_nurses ;
4 int : total_days =10;
5 set of int : DAYS = 1.. total_days ;
6 int : required_days =3;
7 int : required_nights =2;
8 int : minimum_nights =2;
9 array [ DAYS ] of var set of NURSES : timeArrow1 ;
10 array [ DAYS ] of var set of NURSES : timeArrow2 ;
11 constraint forall (i in 1.. total_days )(
12     card ( timeArrow1 [ i ])== required_nights /\
13     card ( timeArrow2 [ i ])== required_days /\
14     disjoint ( timeArrow1 [i], timeArrow2 [i] );
15 constraint forall (i in 3.. total_days )(
16     disjoint ( timeArrow1 [i],
17     timeArrow1 [i -2] intersect timeArrow1 [i -
18     1]));
19 constraint forall (j in 1.. total_nurses )(
20     sum (i in 1.. total_days )(
21     bool2int (j in timeArrow1 [i])) >= minimum_nights
22     );
23 int : total_options = 3;
24 set of int : SHIFTS = 1.. total_options ;
25 int : day_shift = 1;
26 int : night_shift = 2;
27 int : off_shift = 3;
28 array [ SHIFTS ] of string : options = ["d","n","-"];
29 array [NURSES , DAYS ] of var SHIFTS : roster ;
30 constraint forall (j in 1.. total_nurses )(
31     forall (i in 1.. total_days )(
32     roster [j,i] = if j in timeArrow1 [i]
33     then night_shift
34     elseif j in timeArrow2 [i] then
35     day_shift
36     else off_shift endif ));
37 solve satisfy ;
38 output [ options [ fix ( roster [i,j] )]
39     ++ if j== total_days then "\n" else " " endif
40     | i in NURSES , j in DAYS ];

```

4. CONCLUSIONS

En route to reformulate a REGULAR-based into SLIDE-based solution to nurse rostering by simplifying the full probability tree of shifts, the author of this paper

proposes the metaphor of gamut. It may be compared to set-covering formalization that may be contributed to [44] in his elaboration on the possibility of using linear programming in scheduling in a paper entitled "Traffic Delays at Toll Booths" [45]. His formalization assumes a variable for every possible shift.

	Day1	Day2	Day3	Day4	Day5	Day6	Day7	Day8	Day9	Day10
Nurse1					Night			Night		
Nurse2					Night			Night		
Nurse3	Day	Day	Night	Night		Night	Night		Night	Night
Nurse4	Day	Day	Day	Day		Night	Night		Night	Night
Nurse5	Day	Day	Night	Night	Day	Day	Day	Day	Day	Day
Nurse6	Night	Night	Day	Day	Day	Day	Day	Day	Day	Day
Nurse7	Night	Night	Day	Day	Day	Day	Day	Day	Day	Day

Fig. 4. Roster.

A gamut, on the other hand, provides a compact nomenclature for a SET of all possible values for a VECTOR, spanning certain basis dimensions. But, the gamut, as seen in nurse rostering example, can be a bit elusive, as it is constituted from three consecutive shifts and has an exclusion of the state of three consecutive night shifts.

Through elaboration of Minizinc implementation, time arrow metaphor is presented and exploited. Decomposing the time arrow into a two or more arrows, may help in parallelizing the execution by applying constraints on corresponding time arrows.

5. FUTURE WORK

There could be specialized propagators for the knitting and the anticipatory ways to be proposed. Yet, one possible future direction is to attempt generate fair rosters [46]. Another direction is to extend the proposed work of nurse-to-shift assignment, by considering other processes of planning decisions in health care [47].

A philosophical future direction may be the consideration of bewildering aspects of the time arrow such as the reversibility [41], and paradoxes [48].

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